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#### COMP 2211 Exploring Artificial Intelligence Minimax and Alpha-beta Pruning Dr. Desmond Tsoi

Department of Computer Science & Engineering The Hong Kong University of Science and Technology, Hong Kong SAR, China



### **Overview**

- In this topic, you will learn how to implement an AI for a 2-player game, Tic-Tac-Toe.
- The AI agent makes use of searching strategies to optimally play the game.
	- Brute-force† : Minimax
	- Minimax with Alpha-Beta Pruning
- The same AI strategies can also be applied to other 2-player games, such as Connect-four, Chess, Checkers, Backgammon, etc.  $\overline{T_{\text{IC-Tac-Toe}}}$









Your turn.



Connect-four Chess Checkers Backgammon

†Brute force refers to straightforward methods of solving a problem.

#### Games

- Games are the oldest, most well-studied domain in artificial intelligence.
- Why?
	- They are fun.
	- Easy to represent, rules are clear.
	- Possible combination of move can be big, e.g., there are  $\sim 10^{154}$  possible moves in chess.
	- Like the "real world" in that decisions have to be made and time is important.
	- Easy to determine when a program is doing well.



# Types of Game

#### • Perfect vs Imperfect information game

- Perfect: Player knows all the possible moves of himself and opponent and their results.
- Imperfect: Player does not know all the possible moves of the opponent.

#### Zero-sum vs Non-zero-sum game

- Zero-sum: A two person game where one player's gain is equal to the other player's loss on any given play of the game.
- Non-zero-sum: A two person game where player's gain is not necessarily equal to the other player's loss on any given play of the game.

#### Deterministic vs Non-deterministic game

- Deterministic: A game that does not involve random choices.
- Non-deterministic: A game that involves some random choices.

## **Examples**



#### **o** Tic-Tac-Toe

- Has perfect information: All players know the moves previously made by all other players.
- $\bullet$  Is zero-sum, as it involves two players, if player 1 wins, she can score 1 and player 2 who losses score (-1) and if the game ends in draw, both players score 0
- Is deterministic: it is a game that does not involve random choices

#### Question

Could you explain the others? For example, why chess game has perfect information, is zero-sum and deterministic?

#### Tic-Tac-Toe

- Tic-Tac-Toe (also called Noughts and Crosses) is a paper-and-pencil game for two players, who take turns marking the spaces in a  $3\times3$  grid.
- The player who succeeds in placing three of their marks in a diagonal, horizontal, or vertical line is the winner.





- First, if there is a move that will win the game, play it.
- If not, check if there is a move that will block a win from the opposing player. If so, place your marker in the blocking spot.
- Otherwise, just place your marker in a random empty cell.
- The AI may not play optimally, i.e., the AI may not achieve the best results.



# AI - Brute Force Approach

Enumerate all the possible states of the game.



- At the very beginning, a player can choose any of the 9 available empty spaces on the game board.
- We can represent these 9 available moves using a Tree as shown on the left.

## AI - Brute Force Approach

- 1st action: 9 possible choices at the very beginning of the game
- 2nd action:  $9\times8$  (for each 9 possible first choices, we have 8 possible remaining choices)
- 3rd action:  $9\times8\times7$  (for each  $9\times8=72$  previous choices we have 7 possible remaining choices)
- $\bullet$  ...
- Last action:  $9! = 9 \times 8 \times 7 \times \cdots \times 1 = 362880$  choices

Total number of states =  $(9 \times 8 \times 7 \times 6 \times 5 \times 4 \times 3 \times 2 \times 1)$  +  $(9 \times 8 \times 7 \times 6 \times 5 \times 4 \times 3 \times 2)$  +  $(9 \times 8 \times 7 \times 6 \times 5 \times 4 \times 3)$  + · · ·  $(9 \times 8) +$ (9)  $= 986409$ 

# AI - Brute Force Approach

#### Questions

- 1. Is it really to have 986409 possible states in the game? Answer: No. The game can end before there is no more empty space on the board.
- 2. How to determine which move to take based on the current game state? Answer: Use minimax. :)



#### Minimax

- Minimax is a recursive algorithm which is used to choose an optimal move for a player assuming that the other player is also playing optimally.
- In minimax, the two players are called maximizer (denoted as MAX) and minimizer (denoted as MIN), and define a scoring method from the standpoint of the MAX player.
- The maximizer (the AI of Tic-Tac-Toe game) tries to maximize the score while the minimizer (the Human player of Tic-Tac-Toe game) tries to minimize the score of AI.
- It is called minimax because it helps in minimizing that the other players can force us to receive.

#### Minimax

- Steps of the minimax algorithm
	- 1. Construct the complete game tree
	- 2. Apply the utility function to each terminal state.
	- 3. Beginning with the terminal states, determine the utility of the predecessor nodes as follows:
		- Node is a MIN-node: value is the minimum of the children nodes
		- Node is a MAX-node: value is the maximum of the children nodes
	- 4. From the initial state (i.e., root of the game tree), MAX chooses the move that leads to the highest value (minimax decision)

Search the game tree in a Depth-First-Search manner to find the value of the root. Depth-first search is an algorithm for traversing or searching tree data structures. The algorithm starts at the root node and explores as far as possible along each branch before backtracking.

#### Utility function

Numeric values given to terminal states. For example, in Tic-Tac-Toe, win  $= +1$ , loss  $= -1$  and draw  $= 0$ .

#### Minimax Example - Initial



- **•** Terminal states' utility: Max wins  $= 1$ , Max loses  $= -1$ ,  $Draw = 0$
- **•** For maximizier, initialize each node to negative infinity,  $-\infty$
- **•** For minimizier, initialize each node to positive infinity,  $\infty$



- **•** For the leftmost state in level 2, the maximizer updates its value with the maximum of  $-\infty$  and 0, where the value 0 is obtained from its child in level 3.
- So, its value is changed from  $-\infty$  to 0.



- **•** For the leftmost state in level 1, the minimizer updates its value with the minimum of  $\infty$  and 0. where the value 0 is obtained from its left child in level 2.
- So, its value is changed from  $\infty$  to 0.



- **For the leftmost state** in level 1, the minimizer updates its value with the minimum of 0 and -1, where the value -1 is obtained from its right child in level 2.
- So, its value is changed from  $0$  to  $-1$ .



- For the state in level 0, the maximizer updates its value with the maximum of  $-\infty$  and -1, where the value -1 is obtained from its left child in level 1.
- So, its value is changed from  $-\infty$  to -1.



- For the state in level 0, the maximizer updates its value with the maximum of -1 and  $+1$ , where the value  $+1$  is obtained from its middle child in level 1.
- So, its value is changed from  $-1$  to  $+1$ .



- For the state in level 2. the maximizer updates its value with the maximum of  $-\infty$  and  $+1$ , where the value  $+1$  is obtained from its child in level 3.
- So, its value is changed from  $-\infty$  to  $+1$ .



- For the state in level 1. the minimizer updates its value with the minimum of  $\infty$  and  $+1$ , where the value  $+1$  is obtained from its left child in level 2.
- So, its value is changed from  $\infty$  to  $+1$ .



- For the state in level 2. the maximizer updates its value with the maximum of  $-\infty$  and 0, where the value 0 is obtained from its child in level 3.
- So, its value is changed from  $-\infty$  to 0.



- For the state in level 1. the minimizer updates its value with the minimum of  $+1$  and 0, where the value 0 is obtained from its right child in level 2.
- So, its value is changed from  $+1$  to 0.



- For the state in level 0, the maximizer updates its value with the maximum of  $+1$  and 0, where the value 0 is obtained from its right child in level 1.
- So, its value remains unchanged, i.e. 0.

#### Minimax Example - Done



- $\bullet$  In the top position, it is X's turn.
- $\bullet$  If X plays in the top-center, then O can guarantee a win.
- $\bullet$  If X plays in the left center, X will win.
- $\bullet$  If X plays right center, then O can force a draw.
- Therefore X will choose to play in the left center.

#### Minimax Code from Scratch

```
# Draw the current state of the game
def drawboard(board):
   print("Current state: \n\n");
   for i in range (0,9):
        if board[i] == 0: \qquad # If value is 0, it's an empty cell
        print("- ",end = " ");<br>if board[i] == 1:
                                    # If value is 1, it's the mark of AI
           print("X ", end = " ");
        if board[i] == -1: # If value is -1, it's the mark of human
           print("0", end = " ");if (i+1) % 3 == 0: # Move to the next line after printing 3 elements
           print("n");
   print("\n\ranglen"); # Move the next line after printing the game board
```
# Check if there is a player wins. If human player wins, return -1. # If AI player wins, return 1. If draw, return 0. If the game is not finished, return 2. def check game(board):

```
# There are 8 possible winning patterns.
# 0, 1, 2 (1st row), 3, 4, 5 (2nd row), 6, 7, 8 (3rd row)
# 0, 3, 6 (1st col), 1, 4, 7 (2nd col), 2, 5, 8 (3rd col)
# 0, 4, 8 (Diagonal - top-left to bottom-right)
# 2, 4, 6 (Diagonal - top-right to bottom-left)
winning pos = [ [0,1,2], [3,4,5], [6,7,8], [0,3,6], [1,4,7], [2,5,8], [0,4,8], [2,4,6] ]
```

```
# Check all possible winning patterns for each player.
for i in range(0,8):
    # Check whether the cell is non-empty AND the winning positions with the same mark
    if board[winning pos[i][0]] != 0 and\
       board[winning_pos[i][0]] == board[winning_pos[i][1]] and\
       board[winning_pos[i][0]] == board[winning_pos[i][2]]:
      # Fulfiled the winning condition, return the mark of player
      return board[winning_pos[i][2]];
```

```
# When reaching here, it means no player wins yet
# Check whether there is an empty cell (i.e. with 0).
# If so, the game is not finished. Return 2
if any(element == 0 for element in board): return 2
return 0 # Reach here, it means the game draws
```

```
# Obtain the user's input and update the board
def human turn(board):
   pos = input("Enter O's position [1 to 9]: ");
   pos = int(pos); # Obtain user input, i.e., where to put the mark
   # If the cell picked by the human player is non-empty (i.e., non-zero), illegal
   if board[pos-1] != 0:
      print("Illegal Move!")
      exit(0) # Terminate the program
   # When reaching here, the move is legal and we put -1 (put the mark of human player)
```
board $[pos-1] = -1$ 

```
# Perform minimax search
def minimax(board, player):
   global count \# Refer to the global variable count in main<br>
count += 1 \# Make 1 call of minimax already
                               # Make 1 call of minimax already
   result = check_game(board) # Check if any player wins
   if result != 2: # If the result is not 2, the game is finished,
      return result \qquad # return 1 if the winner is AI.
                               # otherwise return -1 (i.e., human is the winner)
   scores = [] # Create an empty list
   for i in range(0,9): # Check all board locations
       if board[i] == 0: # If the cell is empty
          board[i] = player # Try to put the player's mark at cell i+1# Perform minimax for the next player, and append its score to the list
          scores.append(minimax(board,player*-1))
          board[i] = 0 # Undo the trial
   # Return the max score if the player is AI or
   # return the min score if the player is human
   return max(scores) if player == 1 else min(scores)
```

```
# This function makes computer's move based on minimax.
def ai turn(board):
   pos = -1; # Initialize pos to illegal value, -1, here
   max value = -2; # Initialize max value so far to -2,
                     # which is a value smaller than the min value possible
    for i in range(0,9):
        if board[i] == 0: # If the cell is empty
          board[i] = 1 # Try to put X at cell i+1
          score = minimax(board, -1); # Calculate minimax score for the human player hoard \begin{bmatrix} i \\ i \end{bmatrix} = 0 # Undo the trial
                                      # Undo the trial
          if score > max_value: # If we can a better score in the next level,
             max value = score # update the score and pos
              pos = i
    board[pos] = 1 \qquad # Put X at pos, which is the best move
```

```
# Main function
# The broad is represented in a single dimensional list
# Initalize the board to all zeros (i.e., all empty cells)
board = [ 0, 0, 0, 0, 0, 0, 0, 0, 0]print("Computer: X VS. You: O")
first = input("Play first (Y/N) :")
if first == 'Y' or first == 'y': # If first is 'Y'/'y', human wants to play first
  player = 1
else:
  player = 0
for i in range (0,9): # Play the same. 9 turns in total
   if check_game(board) != 2: # If the value returned by check_game is not 2,
      break # the game is finished
   if (i+player) % 2 == 0: \# Take turns to play the game
      count = 0 \# Count the number of minimax callsai_turn(board)
      print("Count:", count) # Print the count
   else:
      drawboard(board)
      human_turn(board)
     Rm 3553, desmond@ust.hk COMP 2211 (Spring 2022) 30 / 53
```

```
result = check_game(board) # Check the game again
if result==0: # If the result is 0, draw
  drawboard(board)
  print("Draw!!!")
if result==1: # If the result is 1, AI wins
  drawboard(board)
print("AI(X) Wins! Human(0) Loses!")<br>if result==-1: # Tf t
                               # If the result is -1, human wins
  drawboard(board)
  print("Human(O) Wins! AI(X) Loses!")
```
# What if MIN does not play optimally?

- Definition of optimal play for MAX assumes MIN plays optimally, i.e., maximizes worst-case outcome for MAX
- If MIN does not play optimally, MAX will do even better.

#### Disadvantages of Minimax

- It has a huge branching factor (the number of children at each node), which makes the process of reaching the goal state slow.
- It can take a long time to evaluate an entire game tree, especially for games such as Go or Chess
- Search and evaluation of unnecessary nodes or branches of the game tree degrades the overall performance and efficiency of the engine.
- Both min and max players have lots of choices to decide from.
- Exploring the entire tree is not possible as there is a restriction of time and space.

Solution: Alpha Beta Pruning

#### **Observations**

- The minimax algorithm provides us with the optimal path as expected. But did we really need to do all that work? No!
- Let's demonstrate it using an example.



#### **Explanations**

- Let's pretend that we are at the middle level stage of finding the minimum value between 6, 5, and 4.
- We have already calculated the value of node 6. So, we are AT LEAST guaranteed that the minimum value between these 3 nodes is  $\leq$  = 6.
- On the level above (top level), we are trying to maximize the values. We've already calculated a value of 7 to the left, something that is guaranteed to be 6 or less will NEVER BE GREATER than 7.
- So, we don't need to go any further with the calculations.
- How about the minimum value between 3, 2, 1? Same!

# Alpha-Beta Pruning

#### Idea

Prevent ourselves from exploring branches of the game tree that are not worth exploring (because they will have no effect on the final outcome).

#### Alpha and Beta

- We define two values  $\alpha$  and  $\beta$  to store information about the potential of a branch for us or the opponent.
	- $\bullet$   $\alpha$ -cutoff: Terminate branch because AI already has a better opportunity elsewhere.
	- $\theta$ -cutoff: Terminate branch because opponent (human) already has better opportunity elsewhere.
- $\bullet$   $\alpha$  is the best (highest) found so far along the path for Max.
- $\bullet$   $\beta$  is the best (lowest) found so far along the path for Min.
- We can prune a branch if  $\alpha > \beta$ .

## Procedure of Alpha-Beta Pruning

- We define two values  $\alpha$  and  $\beta$  to store information about the potential of a branch for us or the opponent.
- The variable  $\alpha$  and  $\beta$  are used to keep track of the current upper and lower bounds.
- Pass current values of  $\alpha$  and  $\beta$  of parent down to child nodes during search.
- Update values of  $\alpha$  and  $\beta$  during search:
	- For MAX, compare  $\alpha$  with the current node value (v). If  $\alpha < v$ ,  $\alpha = v$ , otherwise  $\alpha$  remains unchanged.
	- For MIN compare  $\beta$  with the current node value (v). If  $\beta > v$ ,  $\beta = v$ , otherwise  $\beta$  remains unchanged.
- Prune remaining branches at a node when  $\alpha \geq \beta$ .
- Update values of  $\alpha$  and  $\beta$  of parent using the value returned by the child.
	- If parent node is MAX, compare  $\alpha$  with the value of child (v). If  $\alpha < v$ ,  $\alpha = v$ , otherwise  $\alpha$ remain unchanged.
	- If parent node is MIN, compare  $\beta$  with the value of child (v). If  $\beta > v$ ,  $\beta = v$ , otherwise  $\beta$ remain unchanged.



- At the first step, the Max player will initialize  $\alpha$  to  $-\infty$ , and  $\beta$  to  $\infty$ .
- We check whether  $\alpha \geq \beta$ . It is not and therefore no pruning is done.



- The Max player moves from root node to its left child by passing down  $\alpha = -\infty$  and  $\beta = \infty$  to it.
- At the left node of the middle level, we do the following with the first terminal value 9.
	- $\bullet$  Node value  $=$ Min(node value, 9), i.e. Min $(\infty, 9) = 9$ .
	- $\beta = \text{Min}(\beta, 9)$ , i.e.  $Min(\infty, 9) = 9.$
- We check whether  $\alpha \geq \beta$ . It is not and therefore no pruning is done.



- The Max player moves from root node to its middle child.
- At the left node of the middle level, we do the following with the second terminal value 8.
	- $\bullet$  Node value  $=$ Min(node value, 8), i.e.  $Min(9.8) = 8$ .
	- $\theta = \text{Min}(\beta, 8)$ , i.e.  $Min(9, 8) = 8.$
- We check whether  $\alpha > \beta$ . It is not and therefore no pruning is done.



- The Max player moves from root node to its right child.
- At the left node of the middle level, we do the following with the second terminal value 7.
	- $\bullet$  Node value  $=$ Min(node value, 7), i.e.  $Min(8.7) = 7$ .
	- $\theta = \text{Min}(\beta, 7)$ , i.e.  $Min(8, 7) = 7.$
- We check whether  $\alpha > \beta$ . It is not and therefore no pruning is done.



- Now, the algorithm backtrack to the root, where the value of the middle left node, i.e. 7, is passed to the parent.
- At the root node, we do the following with the returned value 7.
	- $\bullet$  Node value  $=$ Max(node value, 7), i.e. Max $(-\infty, 7) =$ 7
	- $\alpha = \text{Max}(\alpha, 7)$ , i.e.  $Max(-\infty, 7) = 7$
- We check whether  $\alpha \geq \beta$ . It is not and therefore no pruning is done.

# Alpha-Beta Pruning Example **The Max player moves**



- from root node to its middle child by passing down  $\alpha = 7$  and  $\beta = \infty$ to it.
- At the middle node of the middle level, we do the following with the first terminal value 6.
	- Node value  $=$ Min(node value, 6), i.e. Min $(\infty, 6) = 6$ .
	- $\beta = \text{Min}(\beta, 6)$ , i.e.  $Min(\infty, 6) = 6.$
- We check whether  $\alpha > \beta$ . We notice that  $\alpha = 7$  is greater than  $\beta = 6$ , and therefore pruning is done.



- Now, the algorithm backtrack to the root, where the value of the center node of the middle layer, i.e. 6, is passed to the parent.
- At the root node, we do the following with the returned value 6.
	- $\bullet$  Node value  $=$ Max(node value, 6), i.e. Max $(7, 6) = 7$ •  $\alpha = \text{Max}(\alpha, 7)$ , i.e.  $Max(7, 6) = 7$
- We check whether  $\alpha \geq \beta$ . It is not and therefore no pruning is done.

# Alpha-Beta Pruning Example **The Max player moves**



- from root node to its right child by passing down  $\alpha = 7$  and  $\beta = \infty$ to it.
- At the right node of the middle level, we do the following with the first terminal value 3.
	- Node value  $=$ Min(node value, 3), i.e. Min $(\infty,3) = 3$ .
	- $\beta = \text{Min}(\beta, 3)$ , i.e.  $Min(\infty, 3) = 3.$
- We check whether  $\alpha > \beta$ . We notice that  $\alpha = 7$  is greater than  $\beta = 3$ , and therefore pruning is done.



- Now, the algorithm backtrack to the root, where the value of the right node of the middle layer, i.e. 3, is passed to the parent.
- At the root node, we do the following with the returned value 3.
	- $\bullet$  Node value  $=$ Max(node value, 3), i.e.  $Max(7, 3) = 7$ •  $\alpha = \text{Max}(\alpha, 3)$ , i.e.  $Max(7, 3) = 7$
- We check whether  $\alpha \geq \beta$ . It is not and therefore no pruning is done.



• The best value for the root node is 7. The figure on the left shows the nodes which are computed and the nodes which has never computed. Hence the optimal value for the maximizer is 7 for this example.

# Alpha-Beta Pruning Code

Replace the minimax function that we defined above with the following:

```
# Perform minimax search with alpha-beta pruning
def minimax_abp(board, player, alpha = -float('inf'), beta = float('inf')):
   global count \# Refer to the global variable count in main
   count += 1 # Make 1 call of minimax abp already
   result = check_game(board) # Check if any player wins
   if result != 2: # If the result is not 2, the game is finished,
      return result # return 1 if the winner is AI,
                            # otherwise return -1 (i.e., human is the winner)
   if player == 1:
      value = -float('inf')for i in range(0,9): # Check all board locations
         if board[i] == 0: # If the cell is empty
            board[i] = player \# Try to put the player's mark at cell i+1
            # Perform minimax with alpha-beta pruning for the next player
            # and update current value if needed
            value = max(value, minimax_abp(board, player*-1, alpha, beta))
            alpha = max(alpha, value) # Update alpha with max of current alpha and value
            board[i] = 0 # Undo the trial
            if alpha >= beta: break
```

```
else:
  value = float('inf')
  for i in range(0,9): \qquad # Check all board locations
      if board[i] == 0: # If the cell is empty
          board[i] = player \# Try to put the player's mark at cell i+1
          # Perform minimax with alpha-beta pruning for the next player
          # and update current value if needed
          value = min(value, minimax abp(board, player*-1, alpha, beta))beta = min(beta, value) \# Update beta with min of current beta and value board [i] = 0 \# Undo the trial
                                     # Undo the trial
          if alpha >= beta: break
```

```
# Return the max score if the player is AI or
# return the min score if the player is human
return value
```
### Advantages of Alpha-Beta Pruning

- It plays a great role in reducing the number of nodes which are found out by minimax.
- When one chance or option is found at the minimum, it stops assessing a move.
- This method also helps to improve the search procedure in an effective way.

Cat makes the best alarm clock



## Disadvantages of Alpha-Beta Pruning

- Alpha-beta prunes large part of search space, but still needs to search all the way to terminal states.
- Note that for some games, the moves must be made in reasonable amount of time.

#### Approaches to Remedy the Problem

- Set a depth limit
- Heuristic evaluation function: Estimated desirability or utility of position
	- $\Rightarrow$  Correlate with the actual chance of winning.

# Alpha-Beta Pruning

#### Question

Can you find the optimal move for the Tic-Tac-Toe example using Alpha-Beta Pruning? :)



That's all! Any questions?

